

Appl. No. 10/063,771
 Amdt. dated April 20, 2006
 Reply to Office action of January 09, 2006

Amendments to the Claims:

This listing of claims will replace all prior versions and listings of claims in the application:

Listing of Claims:

1 (currently amended): A method for wireless communications comprising:

- 5 providing a multicarrier direct-sequence code-division multiple-access (MS-DS/CDMA) communications system;
 generating a code tree of two-dimensional orthogonal variable spreading factor (2D-OVSF) codes, wherein each node of the code tree has a corresponding matrix;
 10 selecting an $M \times N$ matrix from a node of the code tree, where M relates to the number of available frequency carriers in the MS-DS/CDMA system, N relates to a spreading factor, $M=2^k$, $N=2^{k+\alpha}$, k is greater than zero, and α is greater than or equal to zero; and
 assigning the $M \times N$ matrix to a MS-DS/CDMA-enabled device of the
 15 MS-DS/CDMA system to serve as a signature sequence of the device;
 wherein generating the code tree comprises:

- providing a first set of orthogonal 2×2 matrices $\{A^{(1)}_{(2 \times 2)}, A^{(2)}_{(2 \times 2)}\}$;
 providing a second set of orthogonal 2×2 matrices $\{B^{(1)}_{(2 \times 2)}, B^{(2)}_{(2 \times 2)}\}$;
 utilizing the first set of 2×2 matrices to generate a pair of progenitor nodes in the
 20 code tree that respectively represent matrices $A^{(1)}_{(2 \times 2^a)}$ and $A^{(2)}_{(2 \times 2^a)}$ by iterating the relationship:

$$A^{(1)}_{(2 \times 2^{1+\beta})} = [A^{(1)}_{(2 \times 2^\beta)} \quad A^{(2)}_{(2 \times 2^\beta)}],$$

$$A^{(2)}_{(2 \times 2^{1+\beta})} = [A^{(1)}_{(2 \times 2^\beta)} \quad -A^{(2)}_{(2 \times 2^\beta)}]; \text{ and}$$

- utilizing the matrices $A^{(1)}_{(2 \times 2^a)}$ and $A^{(2)}_{(2 \times 2^a)}$ to generate a child node of one of
 25 the progenitor nodes, the child node having the $M \times N$ matrix, by iterating the relationship:

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$$A^{(i-1)}_{(O \times P)} = [B^{(1)}_{(2 \times 2)} \otimes A^{(i/2)}_{(O/2 \times P/2)}],$$

$$A^{(i)}_{(O \times P)} = [B^{(2)}_{(2 \times 2)} \otimes A^{(i/2)}_{(O/2 \times P/2)}];$$

wherein \otimes is a Kronecker product, and children nodes of the code tree correspond to data transmission rates that are slower than those of parent nodes, thereby enabling multirate transmissions by utilizing orthogonal matrices in the code tree.

2 (cancelled).

10 3 (original): The method of claim 1 wherein children nodes of the progenitor nodes are in the form of a binary tree.

4 (original): The method of claim 1 wherein u is greater than zero.

15 5 (original): The method of claim 1 wherein the first set of orthogonal 2×2 matrices is identical to the second set of orthogonal 2×2 matrices.

6 (original): The method of claim 1 wherein any two matrices at an identical layer in the code tree are orthogonal to each other, and any two matrices at different layers in the code tree are orthogonal if and only if one of the matrices is not the parent code of the other matrix.

7-10 (cancelled).

25 11 (currently amended): A wireless communications device comprising:
 a multicarrier direct-sequence code-division multiple-access (MS-DS/CDMA) module for providing MC-DS/CDMA functionality to the wireless device

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according to a signature sequence; and
 a memory for storing a code tree of two-dimensional orthogonal variable spreading factor (2D-OVSF) codes that are capable of serving as the signature sequence, the 2D-OVSF codes each expressible as an $M \times N$ matrix $A^{(i)}_{(M \times N)}$ in which M relates to the number of available frequency carriers available to the wireless device, N relates to a spreading factor, $M=2^k$, $N=2^{k+\alpha}$, k is greater than zero, i ranges from one to M , and α is greater than or equal to zero, the 2D-OVSF codes of the code tree being interrelated by:

$$\begin{aligned} A^{(1)}_{(2 \times 2^{1+\beta})} &= [A^{(1)}_{(2 \times 2^\beta)} \quad A^{(2)}_{(2 \times 2^\beta)}], \\ A^{(2)}_{(2 \times 2^{1+\beta})} &= [A^{(1)}_{(2 \times 2^\beta)} \quad -A^{(2)}_{(2 \times 2^\beta)}], \\ A^{(i-1)}_{(O \times P)} &= [B^{(1)}_{(2 \times 2)} \otimes A^{(i/2)}_{(O/2 \times P/2)}], \text{ and} \\ A^{(i)}_{(O \times P)} &= [B^{(2)}_{(2 \times 2)} \otimes A^{(i/2)}_{(O/2 \times P/2)}]; \end{aligned}$$

wherein \otimes is a Kronecker product, $A^{(1)}_{(2 \times 2)}$ is orthogonal to $A^{(2)}_{(2 \times 2)}$, and $B^{(1)}_{(2 \times 2)}$ is orthogonal to $B^{(2)}_{(2 \times 2)}$; and

wherein children nodes of the code tree correspond to data transmission rates that are slower than those of parent nodes, thereby enabling multirate transmissions by utilizing orthogonal matrices in the code tree.

12 (original): The wireless device of claim 11 wherein α is greater than zero.

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13 (original): The wireless device of claim 11 wherein $A^{(1)}_{(2 \times 2)}$ equals $B^{(1)}_{(2 \times 2)}$, and $A^{(2)}_{(2 \times 2)}$ equals $B^{(2)}_{(2 \times 2)}$.

14 (currently amended): A method for providing a signature sequence to a mobile unit in a multicarrier direct-sequence code-division multiple-access (MS-DS/CDMA)

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communications system, the method comprising:

a base station generating an $M \times N$ matrix according to a matrix generation method;
 the base station transmitting the $M \times N$ matrix to a mobile device; and
 the mobile device utilizing the $M \times N$ matrix to serve as a signature sequence of the
 5 mobile device;

wherein the matrix generation method comprises:

providing a first set of orthogonal 2×2 matrices $\{A^{(1)}_{(2 \times 2)}, A^{(2)}_{(2 \times 2)}\}$;

providing a second set of orthogonal 2×2 matrices $\{B^{(1)}_{(2 \times 2)}, B^{(2)}_{(2 \times 2)}\}$;

utilizing the first set of 2×2 matrices to generate a pair of progenitor nodes in a code
 10 tree that respectively represent matrices $A^{(1)}_{(2 \times 2^a)}$ and $A^{(2)}_{(2 \times 2^a)}$ by iterating
 the relationship:

$$A^{(1)}_{(2 \times 2^{1+b})} = [A^{(1)}_{(2 \times 2^b)} \quad A^{(2)}_{(2 \times 2^b)}],$$

$$A^{(2)}_{(2 \times 2^{1+b})} = [A^{(1)}_{(2 \times 2^b)} \quad -A^{(2)}_{(2 \times 2^b)}]; \text{ and}$$

utilizing the matrices $A^{(1)}_{(2 \times 2^a)}$ and $A^{(2)}_{(2 \times 2^a)}$ to generate a child node of one of
 15 the progenitor nodes, the child node having the $M \times N$ matrix, by iterating the
 relationship:

$$A^{(i-1)}_{(O \times P)} = [B^{(1)}_{(2 \times 2)} \otimes A^{(i/2)}_{(O/2 \times P/2)}],$$

$$A^{(i)}_{(O \times P)} = [B^{(2)}_{(2 \times 2)} \otimes A^{(i/2)}_{(O/2 \times P/2)}];$$

wherein \otimes is a Kronecker product, and children nodes of the code tree correspond to
 20 data transmission rates that are slower than those of parent nodes, thereby
enabling multirate transmissions by utilizing orthogonal matrices in the code
tree.

15 (cancelled).

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16 (new): The method of claim 14 wherein children nodes of the progenitor nodes are in the form of a binary tree.

17 (original): The method of claim 14 wherein α is greater than zero.

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18 (original): The method of claim 14 wherein the first set of orthogonal 2×2 matrices is identical to the second set of orthogonal 2×2 matrices.

10 19 (original): The method of claim 14 wherein any two matrices at an identical layer in the code tree are orthogonal to each other, and any two matrices at different layers in the code tree are orthogonal if and only if one of the matrices is not the parent code of the other matrix.

20-21 (cancelled).

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